

Advanced Digital Signal Processing

Part 4: DFT and FFT

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Definitions

Basic definitions (assumed to be known from lectures about signals and systems):

The Discrete Fourier Transform (DFT):

$$v(n) \longrightarrow V_M(\mu) = \sum_{n=0}^{M-1} v(n) e^{-j\frac{2\pi}{M}\mu n} = \sum_{n=0}^{M-1} v(n) W_M^{\mu n}.$$

The *inverse Discrete Fourier Transform* (*IDFT*):

$$V_M(\mu) \bullet 0 v(n) = \frac{1}{M} \sum_{\mu=0}^{M-1} V_M(\mu) e^{j\frac{2\pi}{M}\mu n} = \frac{1}{M} \sum_{\mu=0}^{M-1} V_M(\mu) W_M^{-\mu n},$$

with the so-called *twiddle factors*

$$W_M = e^{-j\frac{2\pi}{M}}$$

and ${\cal M}$ being the number of DFT points.





Linear and Circular Convolution – Part 1

Basic definitions of both types of convolutions

A *linear convolution* of two sequences $v_1(n)$ and $v_2(n)$ with $n \in \mathbb{Z}$ is defines as

$$y_{\text{lin}}(n) = v_1(n) * v_2(n) = v_2(n) * v_1(n)$$

= $\sum_{\kappa = -\infty}^{\infty} v_1(\kappa) v_2(n-\kappa) = \sum_{\kappa = -\infty}^{\infty} v_2(\kappa) v_1(n-\kappa).$

A *circular convolution* of two periodic sequences $v_1(n)$ and $v_2(n)$ with $n \in \{0, 1, ..., M_{1,2} - 1\}$ and with the same period $M_1 = M_2 = M$ if defined as

$$y_{cir}(n) = v_1(n) \circledast v_2(n) = v_2(n) \circledast v_1(n)$$
$$= \sum_{\kappa=\kappa_0}^{\kappa_0+M-1} v_1(\kappa) v_2(n-\kappa) = \sum_{\kappa=\kappa_0}^{\kappa_0+M-1} v_2(\kappa) v_1(n-\kappa)$$

The parameter κ_0 needs only to fulfill $\kappa_0 \in \mathbb{Z}$.





Linear and Circular Convolution – Part 2

The DFT and it's relation to circular convolution – Part 1:

The DFT is defined as the transform of the periodic signal with length M. Thus, we have

 $y(n) \circ y(n) \bullet Y_M(\mu) \bullet y(n) = y(n + \lambda M).$

Applying the DFT to a circular convolution leads to

$$y(n) = v_1(n) \circledast v_2(n) \frown V_M(\mu) = V_{1,M}(\mu) \cdot V_{2,M}(\mu),$$

with

$$v_1(n) \quad \frown \quad V_{1,M}(\mu),$$

 $v_2(n) \quad \frown \quad V_{2,M}(\mu).$

This means that a circular convolution can be performed very efficiently (see next slides) in the DFT domain!





Linear and Circular Convolution – Part 3

The DFT and it's relation to circular convolution – Part 2:

Proof of the DFT relation with the circular convolution on the blackboard ...





Linear and Circular Convolution – Part 4

Example:



DFT and linear convolution for finite-length sequences – Part 1

Basic ideas

- □ The *filtering operation* can also be carried out in the *frequency domain* using the DFT. This is very attractive, since fast algorithms (fast Fourier transforms) exist.
- The DFT only realizes a *circular convolution*. However, the desired operation for linear filtering is linear convolution. How can this be achieved by means of the DFT?

Given a finite-length sequence $v_1(n)$ with length M_1 and $v_2(n)$ with length M_2 :

□ The *linear convolution* is defined as:

$$y(n) = \sum_{\kappa=0}^{M_1-1} v_1(\kappa) v_2(n-\kappa),$$

with a length $M_1 + M_2 - 1$ of the convolution result y(n).

□ The *frequency domain equivalent* is

$$Y(e^{j\Omega}) = V_1(e^{j\Omega}) V_2(e^{j\Omega}).$$



DFT and linear convolution for finite-length sequences – Part 2

Given a finite-length sequence $v_1(n)$ with length M_1 and $v_2(n)$ with length M_2 (continued):

- In order to represent the sequence y(n) uniquely in the frequency domain by samples of its spectrum $Y(e^{j\Omega})$, the number of samples must be equal or exceed $M_1 + M_2 1$. Thus, a DFT of size $M \ge M_1 + M_2 1$ is required.
- □ Then, the DFT of the linear convolution

$$y(n) = v_1(n) * v_2(n)$$

is
 $Y_M(\mu) = V_{1,M}(\mu) \cdot V_{2,M}(\mu).$

Explanation on the blackboard (if required) ...

This result can be summarized as follows:

□ The circular convolution of two sequences $v_1(n)$ with length M_1 and $v_2(n)$ with length M_2 leads to the same result as the linear convolution $v_1(n) * v_2(n)$ when the lengths of $v_1(n)$ and $v_2(n)$ are increased to $M_1 + M_2 - 1$ by *zero padding*.



Linear Filtering in the DFT Domain – Part 3

DFT and linear convolution for finite-length sequences – Part 3

Alternative interpretation:

- The circular convolution can be interpreted as a linear convolution with aliasing.
- The inverse DFT leads to the following sequence in the time-domain:

$$y_{p}(n) = \begin{cases} \sum_{\lambda=-\infty}^{\infty} y(n-\lambda M), \\ \text{if } 0 \leq n < M, \\ 0, \\ \text{else.} \end{cases}$$

□ For clarification, see example on the right.







Linear Filtering in the DFT Domain – Part 4

DFT and linear convolution for infinite or long sequences – Part 1

Basic objective:

 \Box Filtering a long input signal v(n) with a finite impulse response h(n) of length M_2 :

$$y(n) = \sum_{\kappa=0}^{M_2-1} h(\kappa) v(n-\kappa).$$

First possible realization: the overlap-add method

□ Segment the input signal into separate (non-overlapping) blocks:

$$v_{\nu}(n) = \begin{cases} v(n - \nu M_1), & \text{if } 0 \le n < M_1, \\ 0, & \text{else.} \end{cases}$$

□ Apply zero-padding for the signal blocks $v_{\nu}(n) \rightarrow \tilde{v}_{\nu}(n)$ and for the impulse response $h(n) \rightarrow \tilde{h}(n)$ to obtain a block length $M \ge M_1 + M_2 - 1$. The non-segmented input signal can be reconstructed according to

$$v(n) = \sum_{\nu=-\infty}^{\infty} \tilde{v}_{\nu}(n-\nu M_1).$$



DFT and linear convolution for infinite or long sequences – Part 2

First possible realization: the overlap-add method (continued)

 \Box Compute *M*-point DFTs of $\tilde{v}_{\nu}(n)$ and $\tilde{h}(n)$ (need to be done only once) and multiply the results:

 $Y_{\nu,M}(\mu) = \widetilde{V}_{\nu,M}(\mu) \widetilde{H}_M(\mu), \text{ with } \mu \in \{0, ..., M-1\}.$

- □ The *M*-point inverse DFT $y_{\nu}(n) = \text{IDFT}\{Y_{\nu,M}(\mu)\}$ yields data blocks that are *free from aliasing due to the zero padding* applied before.
- □ Since each input data block $v_{\nu}(n)$ is terminated with $M M_1$ zeros, the last $M M_1$ signal samples from each output block $y_{\nu}(n)$ must be *overlapped* with (added to) the first $M M_1$ signal samples of the succeeding block (linearity of convolution):

$$y(n) = \sum_{\nu=-\infty}^{\infty} \tilde{y}_{\nu}(n-\nu M_1).$$

As we will see later on, this can result in an immense reduction in computational complexity (compared to the direct time-domain realization) since efficient computations of the DFT and IDFT exist.





Linear Filtering in the DFT Domain – Part 6

DFT and linear convolution for infinite or long sequences – Part 3







Linear Filtering in the DFT Domain – Part 7

DFT and linear convolution for infinite or long sequences – Part 4

Second possible realization: the *overlap-save method*

 \Box Segment the input signal into blocks of length M with an overlap of length $M - M_1$:

$$v_{\nu}(n) = \begin{cases} v(n - \nu M_1), & \text{if } 0 \le n < M, \\ 0, & \text{else.} \end{cases}$$

• Apply zero-padding for the impulse response $h(n) \rightarrow \tilde{h}(n)$ to obtain a block length $M \ge M_1 + M_2 - 1$.

Compute *M*-point DFTs of $v_{\nu}(n)$ and $\tilde{h}(n)$ (need to be done only once) and multiply the results:

$$Y_{\nu,M}(\mu) = V_{\nu,M}(\mu) \widetilde{H}_M(\mu), \text{ with } \mu \in \{0, ..., M-1\}.$$



DFT and linear convolution for infinite or long sequences – Part 5

Second possible realization: the overlap-save method (continued)

- □ The *M*-point inverse DFT $y_{\nu}(n) = \text{IDFT}\{Y_{\nu,M}(\mu)\}$ yields data blocks of length *M* with aliasing in the first $M M_1$ samples. These samples must be *discarded*. The last $M_1 = M - M_2 - 1$ samples of $y_{\nu}(n)$ are exactly the same as the result of a linear convolution.
- □ In order to avoid the loss of samples due to aliasing the last $M M_1$ samples are saved and *appended* at the beginning of the next block. The processing is started by setting the first $M M_1$ samples of the first block to zero.





Linear Filtering in the DFT Domain – Part 9

DFT and linear convolution for infinite or long sequences – Part 6

Second possible realization: the *overlap-save method* (continued)







DFT and linear convolution for infinite or long sequences – Part 7

Partner work – Please think about the following questions and try to find answers (first group discussions, afterwards broad discussion in the whole group).
What are the differences between the overlap-add and the overlap-save method?
Are there advantages and disadvantages?
Can you think of applications where you would prefer either overlap-save or overlap-add? Please explain your choice!





Frequency analysis of stationary signals – Leakage effect – Part 1

Preprocessing for spectral analysis of an analog signal v(t) in practice:

 \Box Anti-aliasing lowpass filtering and sampling with $\omega_s > 2\omega_b$, ω_b denoting the cut-off frequency of the signal.

□ For practical purposes (delay, complexity):

Limitation of the signal duration to the time interval $T_0 = LT$

(L: number of samples under consideration, T: sampling interval).

Consequence of the duration limitation:

□ The limitation to a signal duration of T_0 can be modeled as multiplication of the sampled input signal v(n) with a rectangular *window* w(n):

$$\hat{v}(n) = v(n) \cdot w(n)$$
 with $w(n) = \begin{cases} 1, & \text{if } 0 \le n \le L-1, \\ 0, & \text{else.} \end{cases}$





Linear Filtering in the DFT Domain – Part 12

Frequency analysis of stationary signals – Leakage effect – Part 2

The *consequence of the duration limitation* is shown using an example:

Suppose that the input sequence consists of a single sinusoid

 $v(n) = \cos\left(\Omega_0 n\right).$

The Fourier transform is

$$V(e^{j\Omega}) = \pi \Big[\delta_0(\Omega - \Omega_0) + \delta_0(\Omega + \Omega_0) \Big].$$

The Fourier transform of the window function w(n) can be obtained as

$$W(e^{j\Omega}) = \sum_{n=0}^{L-1} e^{-j\Omega n} = \frac{1 - e^{-j\Omega L}}{1 - e^{-j\Omega}} = e^{-j\Omega(L-1)/2} \cdot \frac{\sin(\Omega L/2)}{\sin(\Omega/2)}.$$

The Fourier transform of the windowed sequence $\hat{v}(n)$ is

$$\hat{V}(e^{j\Omega}) = \frac{1}{2\pi} \Big[V(e^{j\Omega}) \circledast W(e^{j\Omega}) \Big]$$

$$= \frac{1}{2} \Big[W(e^{j(\Omega - \Omega_0)}) + W(e^{j(\Omega + \Omega_0)}) \Big].$$



Linear Filtering in the DFT Domain – Part 13

Frequency analysis of stationary signals – Leakage effect – Part 3

Magnitude frequency response $|\hat{V}(e^{j\Omega})|$ for L = 25 and $\Omega_0 = 0, 2 \cdot 2\pi$:



The windowed spectrum $\hat{V}(e^{j\Omega})$ is not localized to the frequency of the cosine v(n) any more. It is spread out over the whole frequency range.

 \Rightarrow This is called "*spectral leaking*".



Frequency analysis of stationary signals – Leakage effect – Part 4

Properties of the rectangle windowing:

 \Box First zero crossing of $W(e^{j\Omega})$ at $\Omega_z = \pm 2\pi/L$

- $\square \Rightarrow$ The larger the number of sampling points L (and thus also the width of the rectangular window) the smaller becomes Ω_z (and thus the main lobe of the frequency response).
- □ ⇒ Decreasing the frequency resolution (making the window width smaller) leads to an increase of the time resolution and vice versa.
 Duality of time and frequency domain.

Practical scope of the DFT:

Use of the DFT in order to obtain a sampled representation of the spectrum according to

$$\widehat{V}(e^{j\Omega_{\mu}}), \quad \Omega_{\mu} = \frac{2\pi}{M}\mu, \quad \mu = 0, ..., M - 1.$$





Linear Filtering in the DFT Domain – Part 15

Frequency analysis of stationary signals – Leakage effect – Part 5

Special case: If
$$M = L$$
 and $\Omega = \frac{2\pi}{M}\nu$, $\nu = 0, ..., M-1$

then the Fourier transform is exactly zero at the sampled frequencies Ω_{μ} except for $\mu = \nu$.

Example: $M = 50, n = 0, \dots, M - 1$, rectangular window w(n)

$$v_0(n) = \cos\left(1.0\frac{2\pi}{M}n\right), \quad v_1(n) = \cos\left(1.2\frac{2\pi}{M}n\right).$$

Results:

$$\Box \operatorname{DFT} \operatorname{of} \hat{v}_0(n) = v_0(n) w(n) :$$
$$\widehat{V}_0(e^{j\Omega}) = \frac{1}{2\pi} \bigg[V_0(e^{j\Omega}) \circledast W(e^{j\Omega}) \bigg] = 0 \quad \text{for } \Omega_\mu = \mu 2\pi/M$$

except for $\mu = 1$ since Ω_0 is exactly an integer multiple of $2\pi/M$.

 \Rightarrow The periodic repetition of $v_0(n)$ leads to a pure cosine.

□ DFT of
$$v_1(n)$$
: $\Omega_{\nu} \neq \nu 2\pi/M \Rightarrow \widehat{V}(e^{j\Omega_{\mu}}) \neq 0$ for $\Omega_{\mu} = \mu 2\pi/M$
⇒ The periodic repetition of $v_1(n)$ is not a cosine sequence anymore



Linear Filtering in the DFT Domain – Part 16

Frequency analysis of stationary signals – Leakage effect – Part 6

Example (continued):





Linear Filtering in the DFT Domain – Part 16

Frequency analysis of stationary signals – Leakage effect – Part 7

Example (continued):







Linear Filtering in the DFT Domain – Part 16

Frequency analysis of stationary signals – Leakage effect – Part 8

Example (continued):







Linear Filtering in the DFT Domain – Part 16



Example (continued):

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Frequency analysis of stationary signals – Leakage effect – Part 10

Partner work – Please think about the following questions and try to find answers (first group discussions, afterwards broad discussion in the whole group).

□ Why is the spectrum of a signal that you analyze using a DFT "widened" and "smeared" in general?

What can you do in order to minimize the effect?

.....

.....

□ Why is a longer sequence length not always the better choice?





Frequency analysis of stationary signals – Windowing – Part 1

Windowing not only distorts the spectral estimate due to leakage effects, it also influences the spectral resolution.

First example:

Consider a sequence of two frequency components

$$v(n) = \cos(\Omega_1 n) + \cos(\Omega_2 n),$$

$$\hat{v}(n) = v(n) \cdot w(n)$$

$$\hat{V}(e^{j\Omega}) = \frac{1}{2} \left[W(e^{j(\Omega - \Omega_1)}) + W(e^{j(\Omega - \Omega_2)}) + W(e^{j(\Omega + \Omega_1)}) + W(e^{j(\Omega + \Omega_2)}) \right]$$

where $W(e^{j\Omega})$ is the Fourier transform of the rectangular window w(n).





Frequency analysis of stationary signals – Windowing – Part 2

Consideration of three cases for the relation between Ω_1, Ω_2 and L:

- $\Box 2\pi/L < |\Omega_1 \Omega_2|$: The two maxima (the main lobes) for both window spectra $W(e^{j(\Omega \Omega_1)})$ and $W(e^{j(\Omega \Omega_2)})$ can be separated.
- \square $2\pi/L = |\Omega_1 \Omega_2|$: Correct values for the spectral samples, but the main lobes cannot be separated anymore.
- $\square \quad 2\pi/L > |\Omega_1 \Omega_2| : \text{The main lobes of } W(e^{j(\Omega \Omega_1)}) \text{ and } W(e^{j(\Omega \Omega_2)}) \text{ overlap.}$

The ability to resolve spectral lines of different frequencies is limited by the main lobe width, which also depends on the length of the window impulse response *L*.





Linear Filtering in the DFT Domain – Part 20

Frequency analysis of stationary signals – Windowing – Part 3

Second example:

Depicted (on the next slide) are the frequency responses $|\widehat{V}(e^{j\Omega})|$ for

$$\hat{v}(n) = \left[\cos(\Omega_0 n) + \cos(\Omega_1 n) + \cos(\Omega_2 n)\right] w(n)$$

with $\Omega_0 = 0.2\pi, \ \Omega_1 = 0.22\pi$ and $\Omega_2 = 0.6\pi$ for different window lengths L.





Linear Filtering in the DFT Domain – Part 21

Frequency analysis of stationary signals – Windowing – Part 4

Second example (continued):





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Linear Filtering in the DFT Domain – Part 22

Frequency analysis of stationary signals – Windowing – Part 5

Second example (continued):





Frequency analysis of stationary signals – Windowing – Part 6

Approach to reduce leakage: Other window functions with lower side lobes (however, this comes with an increase of the width of the main lobe).

One possible (often used) window: the *Hann window*, defined as

$$w_{\text{\tiny Han}}(n) = \begin{cases} \frac{1}{2} \left[1 - \cos\left(\frac{2\pi}{L-1}n\right) \right], & \text{if } 0 \le n \le L-1, \\ 0, & \text{else.} \end{cases}$$

Magnitude frequency response $|\hat{V}(e^{j\Omega})|$ of the cosine-function $v(n) = \cos(\Omega_0 n)$ after windowing with the *Hann window*:







Linear Filtering in the DFT Domain – Part 24

Frequency analysis of stationary signals – Windowing – Part 7

Spectrum of the signal $v(n) = \cos(\Omega_0 n) + \cos(\Omega_1 n) + \cos(\Omega_2 n)$ after windowing with the *Hann window*:



The reduction of the sidelobes and the reduced resolution compared to the rectangular window can be clearly observed.





Frequency analysis of stationary signals – Windowing – Part 8

Spectrum of the signal $v(n) = \cos(\Omega_0 n) + \cos(\Omega_1 n) + \cos(\Omega_2 n)$ after windowing with the *Hann window* (continued):





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Linear Filtering in the DFT Domain – Part 26

Frequency analysis of stationary signals – Comparison of window sequences – Part 1

In the following we will *compare different window sequences*. Before we start, some *global remarks*:

 \Box All windows will have the *same length*: L = 50 (which means that 51 coefficients are used).

The windows will be *centered around 0*, meaning that they start at index -25 and end at index 25.

□ From time to time the so-called *Dirichlet kernel* will be used. This an abbreviation for the following function:

$$D_k(\Omega) = \frac{1}{2\pi} \sum_{m=-k}^k e^{j\Omega m}$$
$$= \frac{1}{2\pi} \frac{\sin\left(\frac{k+1}{2}\Omega\right)}{\sin\left(\frac{\Omega}{2}\right)}.$$

In the following all *time sequences* are *normalized* to have *maximum amplitude of one*.
 The *magnitude frequency responses* are *normalized* to have *unity gain at frequency zero*.



Johann Peter Gustav Lejeune Dirichlet, German mathematician, 1805 – 1859 (Source: Wikipedia)




Linear Filtering in the DFT Domain – Part 27

Frequency analysis of stationary signals – Comparison of window sequences – Part 2

Some *global remarks* (continued):

- □ Creation of window sequences:
 - □ One can start with basic sequences (sin/cos, Gauss function, polynomials, ...)
 - □ One can combine basic sequences (by means of addition, multiplication, convolution, ...)
- □ Desired shapes of sequences (their corresponding transforms):
 - □ Small main lobe / large attenuation of side lobes (most of the time)
 - □ Sometimes also a certain width of the main lobe (SONAR, RADAR, ...)
 - Addition of shifted (und multiplied) window sequences (overlap-add filterbanks, ...)
 - □ Small aliasing components (if combined with subsampling in filterbanks, ...)

The more window sequences you know, the better ...



Linear Filtering in the DFT Domain – Part 28







Linear Filtering in the DFT Domain – Part 29





Linear Filtering in the DFT Domain – Part 30

Frequency analysis of stationary signals – Comparison of window sequences – Part 4

Some *problems* you already know ...

... Create windows that go down to zero at the edges. A simple idea is to use a window that has a *triangular* shape (instead of a rectangular one).





Linear Filtering in the DFT Domain – Part 31







Linear Filtering in the DFT Domain – Part 32





Linear Filtering in the DFT Domain – Part 33







Linear Filtering in the DFT Domain – Part 34





Linear Filtering in the DFT Domain – Part 35





Linear Filtering in the DFT Domain – Part 36





Linear Filtering in the DFT Domain – Part 37





Linear Filtering in the DFT Domain – Part 38







Linear Filtering in the DFT Domain – Part 39







Linear Filtering in the DFT Domain – Part 40







Linear Filtering in the DFT Domain – Part 41





Linear Filtering in the DFT Domain – Part 42





Linear Filtering in the DFT Domain – Part 43





Linear Filtering in the DFT Domain – Part 44





Linear Filtering in the DFT Domain – Part 45





Linear Filtering in the DFT Domain – Part 46





Linear Filtering in the DFT Domain – Part 47







Linear Filtering in the DFT Domain – Part 48







Linear Filtering in the DFT Domain – Part 49







Linear Filtering in the DFT Domain – Part 50







Linear Filtering in the DFT Domain – Part 51







Linear Filtering in the DFT Domain – Part 52







Linear Filtering in the DFT Domain – Part 53







Linear Filtering in the DFT Domain – Part 54







Linear Filtering in the DFT Domain – Part 55





Linear Filtering in the DFT Domain – Part 56







Linear Filtering in the DFT Domain – Part 57





Linear Filtering in the DFT Domain – Part 58









Linear Filtering in the DFT Domain – Part 59

Frequency analysis of stationary signals – Comparison of window sequences – Part 34

Could be continued virtually endlessly, but for now it should be enough ...





Linear Filtering in the DFT Domain – Part 59

Frequency analysis of stationary signals – Windowing

Partner work – Please think about the following questions and try to find answers (first group discussions, afterwards broad discussion in the whole group).

□ Why do we apply a window function before performing a Fourier transform?

□ How do you select a window function? What prior information might be useful to know before you chose a window function?

Uwhich window would you chose if you need a narrow main lobe? Is your choice optimal in some sense?





Fast Computation of the DFT: The FFT – Part 1

Basics – Part 1

When computing the *complexity of a DFT* for the complex signals $v(n) \in \mathbb{C}$ and complex spectra $V_M(\mu) \in \mathbb{C}$, we obtain

 $V_M(\mu) = \underbrace{\sum_{n=0}^{M-1} \underbrace{v(n) \cdot e^{-j\mu \frac{2\pi}{M}n}}_{1 \text{ compl. mult.}}}_{M \text{ compl. mult.}} \quad \text{with } \underbrace{\mu \in \{0, 1, ..., M-1\}}_{M \text{ result values}}.$

In total we need about M^2 complex multiplications and additions

Remarks (part 1):

 \Box 1 complex multiplication \rightarrow 4 real multiplications + 2 real additions,

1 complex addition \rightarrow 2 real additions.

□ When looking closer we see that not all operations require the above mentioned complexity:

 \Box *M* values have to be added, \Rightarrow (*M* - 1) additions,

 \Box for the factors e^{j0} , $e^{j\pi\lambda}$, $e^{\pm j\lambda\pi/2}$ multiplications are not required,

 \Box for $\mu = 0$ multiplications are not necessary.





Fast Computation of the DFT: The FFT – Part 2

Basics – Part 2

Remarks (part 2):

Multiplications and additions are not the only operations that should be considered for analyzing the computational complexity.

□ *Memory access operations, checking conditions,* etc. are as important as additions and multiplications.

□ Cost functions for *complexity measures* should be *adapted* to the *individually used hardware*.

Basic idea for reducing the computational complexity:

The basic idea for reducing the complexity of a DFT is to decompose the *"big problem"* into several *"smaller problems"*. This usually leads to a reduction in complexity. However, this *"trick"* can not always be applied.




Fast Computation of the DFT: The FFT – Insertion Part 1

The Goertzel Algorithm

Basics:

- □ The *name* stems for an American physicist (1919 2002).
- The basic idea of this algorithm is to reduce the number of multiplications as much as possible if only a single DFT pin should be computed.
- □ The principle works also if an *arbitrary frequency supporting point* (not on the DFT grid) should be computed.

Number of real multiplication for a single bin:

 \Box The signal v(n) is assumed to be real.

$$V_M(\mu_0) = \underbrace{\sum_{n=0}^{M-1} \underbrace{v(n) \cdot e^{-j\mu_0 \frac{2\pi}{M}n}}_{2 \text{ real mult.}}}_{2M \text{ real mult.}}$$





Fast Computation of the DFT: The FFT – Insertion Part 2

The Goertzel Algorithm

Derivation:

$$\begin{split} V_{M}(\mu_{0}) &= \sum_{n=0}^{M-1} v(n) \cdot e^{-j\mu_{0}\frac{2\pi}{M}n} \\ & \dots \text{ inserting the abbreviation } W_{M} = e^{-j\frac{2\pi}{M}} \text{ (twiddle factor)} \dots \\ &= \sum_{n=0}^{M-1} v(n) \cdot W_{M}^{\mu_{0}n} \\ & \dots \text{ exploiting that } W_{M}^{-\mu_{0}M} = e^{j\frac{2\pi}{M}\mu_{0}M} = e^{j2\pi\mu_{0}} = 1 \dots \\ &= \sum_{n=0}^{M-1} v(n) \cdot W_{M}^{-\mu_{0}(M-n)}. \end{split}$$

Interpretation as a convolution (but only at one specific sample):

$$V_M(\mu_0) = \sum_{n=0}^{M-1} v(n) \cdot W_M^{-\mu_0(M-n)} = \left(v(n) * W_M^{-n\mu_0} \right) \Big|_{n=M}$$





Fast Computation of the DFT: The FFT – Insertion Part 3

The Goertzel Algorithm

Interpretation as a convolution (but only at one specific sample):

$$V_{M}(\mu_{0}) = \sum_{n=0}^{M-1} v(n) \cdot W_{M}^{-\mu_{0}(M-n)} = \left(v(n) * W_{M}^{-(n+1)\mu_{0}} \right) \Big|_{n=M-1}.$$

FIR filter with the coefficients $h_{i} = W_{M}^{-(i+1)\mu_{0}}$ with $i \in \{0, 1, ..., M-1\}$.

Example for n = 4:

v(4) v	v(3)	v(2)	v(1)	v(0)	0	0	0
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$W_M^{-1\mu_0} W_M^{-2\mu_0} W_M^{-3\mu_0}$					$W_M^{-M\mu_0}$
---	--	--	--	--	-----------------

Example for n = M - 1:

v(M-1)v(M-2)v(M-3)					v(0)
--------------------	--	--	--	--	------

$W_M^{-1\mu_0} W_M^{-2\mu_0} W_M^{-3\mu_0}$				••••	$W_M^{-M\mu_0}$
---	--	--	--	------	-----------------





Fast Computation of the DFT: The FFT – Insertion Part 4

The Goertzel Algorithm

Extending the filter:

Example for n = M - 1, "*normal*" filter:

v(M-1)v(M-2)v(M-3)					v(0)	
--------------------	--	--	--	--	------	--

$W_M^{-1\mu_0} W_M^{-2\mu_0}$	$W_M^{-3\mu_0}$					$W_M^{-M\mu_0}$
-------------------------------	-----------------	--	--	--	--	-----------------

Example for n = M - 1, "*extended*" filter:

v(M-1)v(M-2)v(M-3)					v(0)	0	0	
--------------------	--	--	--	--	------	---	---	--





Fast Computation of the DFT: The FFT – Insertion Part 5

The Goertzel Algorithm

Some more specific look on the filter realization:

$$\begin{split} H(z) &= \sum_{i=0}^{\infty} h_i z^{-i} \\ & \text{... inserting the definition of the filter coefficients } h_i = W_M^{-(i+1)\mu_0} \text{...} \\ &= \sum_{i=0}^{\infty} W_M^{-(i-1)\mu_0} z^{-i} \\ & \text{... exclude one of the twiddle factors ...} \\ &= W_M^{\mu_0} \cdot \sum_{i=0}^{\infty} W_M^{-i\mu_0} z^{-i} \\ & \text{... multiply with a "complicatedly written 1"...} \\ &= W_M^{\mu_0} \cdot \frac{1 - W_M^{-\mu_0} z^{-1}}{1 - W_M^{-\mu_0} z^{-1}} \cdot \sum_{i=0}^{\infty} W_M^{-i\mu_0} z^{-i} \end{split}$$





Fast Computation of the DFT: The FFT – Insertion Part 6

The Goertzel Algorithm

Some more specific look on the filter realization (continued):

$$H(z) = W_M^{\mu_0} \cdot \frac{1 - W_M^{-\mu_0} z^{-1}}{1 - W_M^{-\mu_0} z^{-1}} \cdot \sum_{i=0}^{\infty} W_M^{-i\mu_0} z^{-i}$$

... multiply the numerator with the sum ...

$$= W_M^{\mu_0} \cdot \frac{\sum_{i=0}^{\infty} W_M^{-i\mu_0} z^{-i} - \sum_{i=0}^{\infty} W_M^{-(i+1)\mu_0} z^{-(i+1)}}{1 - W_M^{-\mu_0} z^{-1}}$$

... exploit that only the sums differ only in one element ...

$$= W_{M}^{\mu_{0}} \cdot \frac{W_{M}^{0} z^{0}}{1 - W_{M}^{-\mu_{0}} z^{-1}}$$

... simplify the numerator ...
$$= W_{M}^{\mu_{0}} \cdot \frac{1}{1 - W_{M}^{-\mu_{0}} z^{-1}}$$





Fast Computation of the DFT: The FFT – Insertion Part 7

The Goertzel Algorithm

Some more specific look on the filter realization (continued):

$$\begin{split} H(z) &= W_M^{\mu_0} \cdot \frac{1}{1 - W_M^{-\mu_0} z^{-1}} \\ & \dots \text{ multiply with a "complicatedly written 1"...} \\ &= W_M^{\mu_0} \cdot \frac{1 - W_M^{\mu_0} z^{-1}}{1 - W_M^{\mu_0} z^{-1}} \cdot \frac{1}{1 - W_M^{-\mu_0} z^{-1}} \\ & \dots \text{ combine both denominators ...} \\ &= W_M^{\mu_0} \cdot \frac{1 - W_M^{\mu_0} z^{-1}}{1 - 2\cos(\frac{2\pi}{M}\mu_0) z^{-1} + z^{-2}} \\ & \dots \text{ split the filter into three subfilter ...} \\ &= \frac{1}{1 - 2\cos(\frac{2\pi}{M}\mu_0) z^{-1} + z^{-2}} \cdot \left(1 - W_M^{\mu_0} z^{-1}\right) \cdot W_M^{\mu_0} \end{split}$$





Fast Computation of the DFT: The FFT – Insertion Part 8

The Goertzel Algorithm

Some more specific look on the filter realization (continued):

$$H(z) = \underbrace{\frac{1}{1 - 2\cos(\frac{2\pi}{M}\mu_0) z^{-1} + z^{-2}}}_{H_0(z)} \cdot \underbrace{\left(1 - W_M^{\mu_0} z^{-1}\right)}_{H_1(z)} \cdot \underbrace{W_M^{\mu_0}}_{H_2(z)}$$

First filter, computed from n = 0 to n = M - 1:

$$H_0(z) = \frac{1}{1 - 2 \cos(\frac{2\pi}{M}\mu_0)} z^{-1} + z^{-2}}$$

$$y_0(z) = v(n) + b_1 y_0(n-1) - y_0(n-2)$$

Second filter, computed only for n = M - 1:

$$H_1(z) = \left(1 - \underbrace{W_M^{\mu_0}}_{a_1} z^{-1}\right) \qquad \qquad y_1(z) = y_0(n) - a_1 y_0(n-1)$$





Fast Computation of the DFT: The FFT – Insertion Part 9

The Goertzel Algorithm

Some more specific look on the filter realization (continued):

$$H(z) = \underbrace{\frac{1}{1 - 2\cos(\frac{2\pi}{M}\mu_0) z^{-1} + z^{-2}}}_{H_0(z)} \cdot \underbrace{\left(1 - W_M^{\mu_0} z^{-1}\right)}_{H_1(z)} \cdot \underbrace{W_M^{\mu_0}}_{H_2(z)}$$

Third filter, computed only for n = M - 1:





Fast Computation of the DFT: The FFT – Insertion Part 10

The Goertzel Algorithm

Some more specific look on the filter realization (continued):

$$H(z) = \underbrace{\frac{1}{1 - 2\cos(\frac{2\pi}{M}\mu_0) z^{-1} + z^{-2}}}_{H_0(z)} \cdot \underbrace{\left(1 - W_M^{\mu_0} z^{-1}\right)}_{H_1(z)} \cdot \underbrace{W_M^{\mu_0}}_{H_2(z)}$$

The *first filter* requires one real multiplication per sample, in total N real multiplications.

The *second filter* requires two real multiplication (only for the final sample).

The *third filter* requires two real multiplication (only for the final sample).





Fast Computation of the DFT: The FFT – Part 3

Radix-2 approach (decimation in time) – Part 1

For fast (efficient) realizations of the DFT some properties of the so-called twiddle factors $W_M^{\mu n} = e^{-j\mu \frac{2\pi}{M}n}$ can be used. In particular we can utilize

□ the conjugate complex *symmetry*

 $W_M^{-\mu n} = \left[W_M^{\mu n} \right]^*$

 $\hfill\square$ and the *periodicity* both for n and for μ

 $W_M^{\mu n} = W_M^{\mu (n+\lambda M)} = W_M^{(\mu+\lambda M)n}.$

For a so-called *radix 2 realization* of the DFT we *decompose* the input series *into two series of half length*:

$$v_1(n) = v(2n),$$

 $v_2(n) = v(2n+1),$
each with $n \in \{0, 1, ..., \frac{M}{2} - 1\}.$





Fast Computation of the DFT: The FFT – Part 4

Radix-2 approach (decimation in time) – Part 2

If we assume that the orignal series has an *even length*, we can decompose it according to

$$V_{M}(\mu) = \sum_{n=0}^{M-1} v(n) \underbrace{e^{-j\frac{2\pi}{M}\mu n}}_{e^{-j\frac{2\pi}{M}\mu n}}, \quad \mu \in \{0, 1, ..., M-1\}$$

... inserting the definition of the twiddle factors ...

$$= \sum_{n=0}^{M-1} v(n) W_{M}^{\mu n}$$

... splitting the sum into even and odd indices ...

$$= \sum_{n=0}^{M/2-1} \underbrace{v(2n)}_{v_{1}(n)} W_{M}^{2\mu n} + \sum_{n=0}^{M/2-1} \underbrace{v(2n+1)}_{v_{2}(n)} W_{M}^{\mu(2n+1)}$$

... inserting the (signal) definitions from the last slide ...

$$= \sum_{n=0}^{M/2-1} v_{1}(n) \underbrace{W_{M}^{2\mu n}}_{W_{M/2}^{2\mu n}} + W_{M}^{\mu} \sum_{n=0}^{M/2-1} v_{2}(n) \underbrace{W_{M}^{2\mu n}}_{W_{M/2}^{2\mu n}}$$



Fast Computation of the DFT: The FFT – Part 5

Radix-2 approach (decimation in time) – Part 3

DFT decomposition (continued)

$$V_{M}(\mu) = \sum_{n=0}^{M/2-1} v_{1}(n) \underbrace{W_{M}^{2\mu n}}_{W_{M/2}^{\mu n}} + W_{M}^{\mu} \sum_{n=0}^{M/2-1} v_{2}(n) \underbrace{W_{M}^{2\mu n}}_{W_{M/2}^{\mu n}}$$

... inserting the definition of a DFT (of half length) ...

$$= V_{M/2,1}(\mu) + W_{M}^{\mu} V_{M/2,2}(\mu).$$
DFT of length $\frac{M}{2}$ for the signal $v_{2}(n)$
DFT of length $\frac{M}{2}$ for the signal $v_{1}(n)$

Please note that this decomposition is – due to the periodicity of $V_{M/2,1}(\mu)$ and $V_{M/2,2}(\mu)$ – also true for $\mu \geq M/2$.





Fast Computation of the DFT: The FFT – Part 6

Radix-2 approach (decimation in time) – Part 4

For the *computational complexity* we obtain:

Before the decomposition:

 $\longrightarrow M^2$ operations. 1 DFT of order M

□ After the decomposition:

2 DFTs of order
$$M/2 \longrightarrow 2\left(\frac{M}{2}\right)^2 = \frac{2M^2}{4} = \frac{M^2}{2}$$
 operations and

combining the results M operations. \longrightarrow

> Using this "splitting" operation we were able to reduce the complexity form M^2 down to $\frac{1}{2}M^2 + M$. For large orders M this is about a half.





Fast Computation of the DFT: The FFT – Part 7

Radix-2 approach (decimation in time) – Part 5

Graphical explanation of (the first stage of) the decomposition for M = 8:





Fast Computation of the DFT: The FFT – Part 8

Radix-2 approach (decimation in time) – Part 6

This *principle* can be *applied again*. Therefore, M/2 has to be even again. Then we get four DFTs of length M/4 and another M operations for combining those four DFTs such that we get the desired outputs. Together with the operations necessary for combining the low order DFTs we obtain a *complexity* of

$$\frac{M^2}{4} + 2M$$
 complex operations.

We can apply this principle until we reach a "minimum order" DFT of length 2. This can be achieved if we have

 $M = 2^m$, with $m \in \mathbb{N}$.

As a result M has to be a **power of two**.





Fast Computation of the DFT: The FFT – Part 9

Radix-2 approach (decimation in time) – Part 7

Graphical explanation of (the second stage of) the decomposition for M = 8:







Fast Computation of the DFT: The FFT – Part 10

Radix-2 approach (decimation in time) – Part 8

As a last step we have to compute a DFT of length 2. This is achieved by:

 $\tilde{V}_{2}(0) = \tilde{v}(0) + W_{2}^{0} \tilde{v}(1) = \tilde{v}(0) + \tilde{v}(1),$ $\tilde{V}_{2}(1) = \tilde{v}(0) + \tilde{v}(1) = \tilde{v}(0) - \tilde{v}(1),$

 $\tilde{V}_2(1) = \tilde{v}(0) + W_2^1 \tilde{v}(1) = \tilde{v}(0) - \tilde{v}(1).$

As we can see, also over here we need the same *basic scheme* that we have used also in the previous decompositions:



This basic scheme is called *"butterfly"* of a radix-2 FFT. The abbreviation *FFT* stands for *Fast Fourier Transform*.





Fast Computation of the DFT: The FFT – Part 11

Radix-2 approach (decimation in time) – Part 9

When computing the individual butterfly operations we can exploit that the twiddle factors W_M^{ρ} and $-W_M^{\rho}$ differ only in terms of their sign. Thus, we can apply the following simplification:



Pathes without variables using a factor of 1!

This leads to a *further reduction in terms of multiplications* (only 50 % of them are really required).

In total we were able to reduce the required operations for computing a DFT from M^2 down to $M \log_2 M$ by using efficient radix-2 approaches.

Examples:
$$M = 32 \rightarrow M^2 \approx 1000, M \log_2 M \approx 160 \rightarrow$$
 Factor 6,
 $M = 1024 \rightarrow M^2 \approx 10^6, M \log_2 M \approx 10^4 \rightarrow$ Factor 100.





Fast Computation of the DFT: The FFT – Part 12

Radix-2 approach (decimation in time) – Part 10

Graphical explanation of the decomposition for M = 8 with optimized butterfly structure:



Please keep in mind that in each stage only "in-place operations" are required. This means that no now memory has to be allocated for a new stage!





Fast Computation of the DFT: The FFT – Part 13

Radix-2 approach (decimation in time) – Part 11

Graphical explanation of the decomposition for M = 8 (with keeping the "orientation" of the input vector):







Fast Computation of the DFT: The FFT – Part 14

Radix-2 approach (decimation in time) – Part 12

Graphical explanation of the complexity reduction on the black board ...





Fast Computation of the DFT: The FFT – Part 15

Radix-2-decimation-in-time FFT algorithms – Part 13

In-place computations

The intermediate results $V_M^{(\ell)}(\mu_{1,2})$ in the *l*-th stage, l = 0, ..., m-1, are obtained as

 $V_M^{(\ell)}(\mu_1) = V_M^{(\ell-1)}(\mu_1) + W_M^{\rho} V_M^{(\ell-1)}(\mu_2),$ $V_M^{(\ell)}(\mu_2) = V_M^{(\ell-1)}(\mu_1) - W_M^{\rho} V_M^{(\ell-1)}(\mu_2)$

(butterfly computations) where $\mu_1, \mu_2, \rho \in \{0, \dots, M-1\}$ vary from stage to stage.

- Only *M* storage cells are needed, which first contain the values v(n), then the results form the individual stages and finally the values $V_M(\mu)$.
 - \Rightarrow In-place algorithm





Fast Computation of the DFT: The FFT – Part 16

Radix-2-decimation-in-time FFT algorithms – Part 14

Bit-reversal

 \Box The v(n)-values are ordered at the input of the decimation-in-time flow graph in permuted order.

 \Box Example for M = 8, where the indices are written in binary notation:

# flow graph input	000	001	010	011	
v(n)	v(000)	v(100)	v(010)	v(110)	
# flow graph input	100	101	110	111	
v(n)	v(001)	v(101)	v(011)	v(111)	

 \Rightarrow Input data is stored in *bit-reversed* order.

Mirrored at the "center bit" in terms of binary counting!





Fast Computation of the DFT: The FFT – Part 17

Radix-2-decimation-in-time FFT algorithms – Part 15

Bit-reversal

Bit-reversed order is due to the sorting in the even and odd indices in every stage, and thus is also necessary for in-place computation.



[v
ightarrow x, Oppenheim, Schafer, 1999]



Fast Computation of the DFT: The FFT – Part 18

Radix-2-decimation-in-time FFT algorithms – Part 16

Inverse FFT

The inverse DFT is defined as

$$v(n) = \frac{1}{M} \sum_{\mu=0}^{M-1} V_M(\mu) W_M^{-\mu n},$$

that is

$$v(-n) = \frac{1}{M} \sum_{\mu=0}^{M-1} V_M(\mu) W_M^{\mu n}, \iff$$
$$v(M-n) = \frac{1}{M} \operatorname{DFT}\{V_M(\mu)\}$$

 \Rightarrow With additional scaling and index permutations the IDFT can be calculated with the same FFT algorithm as the DFT.





Fast Computation of the DFT: The FFT – Part 19

FFT alternatives – Part 1

Alternative decimation-in-time (DIT) structures

Rearranging of the nodes in the signal flow graphs lead to FFTs with almost arbitrary permutation of the input and output sequence. Reasonable approaches are structures:

- (a) without bit-reversal, or
- (b) bit-reversal in the frequency domain.

The approach (a) has the disadvantage that it is a *non-inplace* algorithm, because the butterfly-structure does not continue past the first stage.

Next slides: the flow graphs for both approaches.





Fast Computation of the DFT: The FFT – Part 20

FFT alternatives – Part 2

Alternative decimation-in-time (DIT) structures (continued)

(a) Flow graph for $M = 8 \ (M \to N, v \to x, V \to X)$



[Oppenheim, Schafer, 1999]





Fast Computation of the DFT: The FFT – Part 21

FFT alternatives – Part 3

Alternative decimation-in-time (DIT) structures (continued)

(b) Flow graph for $M = 8 \ (M \to N, v \to x, V \to X)$



[Oppenheim, Schafer, 1999]





Fast Computation of the DFT: The FFT – Part 22

FFT alternatives – Part 4 Decimation-in-frequency algorithms

Instead of applying the decomposition to time domain

ightarrow Starting the decomposition in the frequency domain

 \rightarrow The sequence of DFT coefficients $V_M(\mu)$ is decomposed into smaller sequences.

 \Rightarrow **Decimation-in-frequency** (DIF) FFT.





Fast Computation of the DFT: The FFT – Part 23

FFT alternatives – Part 5

Decimation-in-frequency algorithms (continued)

Signal flow graph for $M = 8 \ (v \to x, V \to X)$



[Proakis, Manolakis, 1996]





Fast Computation of the DFT: The FFT – Part 24

FFT alternatives – Part 6 Radix r and mixed-radix FFTs

When we generally use

 $M = r^m \quad \text{for } r \ge 2, \quad r, m \in \mathbb{N}$

we obtain DIF or DIT decompositions with a radix r.

Besides r = 2, r = 3 and r = 4 are commonly used.





Fast Computation of the DFT: The FFT – Part 25

FFT alternatives – Part 7

Radix r and mixed-radix FFTs (continued)

Radix-4 butterfly $(q = 0, \dots, M/4 - 1; M \rightarrow N)$:



[Proakis, Manolakis, 1999]





Fast Computation of the DFT: The FFT – Part 26

Convolution of a finite and an infinite sequence:

We will compare now the direct convolution in the time domain with it's DFT/FFT counterpart. For the DFT/FFT realization we will use the overlap-add technique.

For that purpose we will modify a music signal by means of amplifying the low and the very high frequencies of a music recording.







Fast Computation of the DFT: The FFT – Part 27

DFT and FFT:

Partner work – Please think about the following questions and try to find answers (first group discussions, afterwards broad discussion in the whole group).

□ What does "in-place" means and why is this property important for efficient algorithm?

.....

U Which operations should be counted besides multiplications and additions when analyzing the efficiency of an algorithm?

□ How would you realize an FFT of order 180?





Transformation of Real-Valued Sequences – Part 1

If $v(n) \in \mathbb{R} \Rightarrow$ FFT calculation for $v(n) = v_{\rm re}(n) + j v_{\rm im}(n) \in \mathbb{C}$ with $v_{\rm im}(n) = 0$.

 \Rightarrow inefficient due to arithmetic calculation with zero values

In the following: *Methods for efficient usage of a complex FFT for real-valued data*.

DFT of two real sequences – Part 1

Given: $v_1(n), v_2(n) \in \mathbb{R}$ with n = 0, ..., M - 1*Question*: How can we efficiently obtain $V_{1,M}(\mu) = DFT\{v_1(n)\}$ and $V_{2,M}(\mu) = DFT\{v_2(n)\}$ by using the complex DFT?

Define

$$v(n) = v_1(n) + jv_2(n)$$

leading to the DFT

$$V_M(\mu) = \mathrm{DFT}\{v(n)\} = \underbrace{V_{1,M}(\mu)}_{\in \mathbb{C}} + j \underbrace{V_{2,M}(\mu)}_{\in \mathbb{C}}.$$




Transformation of Real-Valued Sequences – Part 2

DFT of two real sequences – Part 2

```
How to separate V_M(\mu) into V_{1,M}(\mu) and V_{2,M}(\mu) ?
```

Symmetry relations of the DFT:

$$v(n) = \underbrace{v_{\text{re,even}}(n) + v_{\text{re,odd}}(n)}_{= v_1(n)} + j \left[\underbrace{v_{\text{im,even}}(n) + v_{\text{im,odd}}(n)}_{= v_2(n)}\right]$$

with the subscripts $\ even\ \mbox{denoting}$ the even part and $\ {\rm odd}$ the odd part.

Corresponding DFTs:

$$\begin{aligned} v_{\text{re,even}}(n) &= \text{IDFT}\{V_{M,\text{re,even}}(\mu)\}, \\ v_{\text{re,odd}}(n) &= j \text{IDFT}\{V_{M,\text{im,odd}}(\mu)\}, \\ j v_{\text{im,even}}(n) &= j \text{IDFT}\{V_{M,\text{im,even}}(\mu)\}, \\ j v_{\text{im,odd}}(n) &= \text{IDFT}\{V_{M,\text{re,odd}}(\mu)\}. \end{aligned}$$





Transformation of Real-Valued Sequences – Part 3

DFT of two real sequences – Part 3

Repetition – symmetry scheme of Fourier transform:



... hope you remember





Transformation of Real-Valued Sequences – Part 4

DFT of two real sequences – Part 4

Thus, we have

$$V_{1,M}(\mu) = \frac{1}{2} \left[V_{M,\text{re}}(\mu) + V_{M,\text{re}}(M-\mu) \right] + \frac{j}{2} \left[V_{M,\text{im}}(\mu) - V_{M,\text{im}}(M-\mu) \right],$$

with

$$V_{M,\text{re,even}}(\mu) = \frac{1}{2} \left[V_{M,\text{re}}(\mu) + V_{M,\text{re}}(M-\mu) \right],$$

$$V_{M,\text{im,odd}}(\mu) = \frac{1}{2} \left[V_{M,\text{im}}(\mu) - V_{M,\text{im}}(M-\mu) \right].$$

Likewise, we have for $V_{2,M}(\mu)$ the relation

$$V_{2,M}(\mu) = \frac{1}{2} \left[V_{M,\text{im}}(\mu) + V_{M,\text{im}}(M-\mu) \right] + \frac{j}{2} \left[V_{M,\text{re}}(\mu) - V_{M,\text{re}}(M-\mu) \right],$$

with

$$V_{M,\text{im,even}}(\mu) = \frac{1}{2} \left[V_{M,\text{im}}(\mu) + V_{M,\text{im}}(M-\mu) \right],$$

$$V_{M,\text{re,odd}}(\mu) = \frac{1}{2} \left[V_{M,\text{re}}(\mu) - V_{M,\text{re}}(M-\mu) \right].$$





Transformation of Real-Valued Sequences – Part 5

DFT of two real sequences – Part 5

Rearranging both relations finally yields

$$V_{1,M}(\mu) = \frac{1}{2} \left[V_M(\mu) + V_M^*(M-\mu) \right],$$

$$V_{2,M}(\mu) = -\frac{j}{2} \left[V_M(\mu) - V_M^*(M-\mu) \right].$$

Due to the hermitean symmetry for real-valued sequences

$$V_{i,M}(\mu) = V_{i,M}^*(M-\mu)$$

only the values for $\mu \in \{0, ..., M/2\}$ have to be calculated.

The values for $\mu \in \{M/2 + 1, ..., M - 1\}$ can be derived from the values for $\mu \in \{0, ..., M/2\}$.

Application:

Fast convolution of two real-valued sequences with the DFT/FFT!





Transformation of Real-Valued Sequences – Part 6

DFT of a 2M-point real sequence – Part 1

Given:
$$v(n) \in \mathbb{R}, n = 0, ..., 2M - 1$$

Wanted:

$$V_{2M}(\mu) = \text{DFT}\{v(n)\} = \sum_{n=0}^{2M-1} v(n) W_{2M}^{\mu n}$$

with $\mu = 0, ..., 2M - 1$.

Hermitian symmetry since $v(n) \in \mathbb{R}$ for all n:

$$V_{2M}(2M-\mu) = V_{2M}^*(\mu), \quad \mu = 0, \dots, M$$

Define

$$\tilde{v}(n) = v(2n) + jv(2n+1), \quad n = 0, ..., M-1,$$

= $v_1(n) + jv_2(n),$

where the even and odd samples of v(n) are written alternatively into the real and imaginary part of $\tilde{v}(n)$.





Transformation of Real-Valued Sequences – Part 7

DFT of a 2M-point real sequence – Part 2

We have a complex sequence $\tilde{v}(n)$ consisting of two real-valued sequences $v_1(n)$ and $v_2(n)$ of length M with the DFT

$$\widetilde{V}_M(\mu') = V_{1,M}(\mu') + j V_{2,M}(\mu'), \quad \mu' = 0, ..., M - 1.$$

 $V_{1,M}(\mu')$ and $V_{2,M}(\mu')$ can easily be obtained as

$$V_{1,M}(\mu') = \frac{1}{2} \left[\widetilde{V}_M(\mu') + \widetilde{V}_M^*(M - \mu') \right],$$

$$V_{2,M}(\mu') = -\frac{j}{2} \left[\widetilde{V}_M(\mu') - \widetilde{V}_M^*(M - \mu') \right]$$

for $\mu' = 0, ..., M - 1$.

In order to calculate $V_{2M}(\mu)$ from $V_{1,M}(\mu')$ and $V_{2,M}(\mu')$ we rearrange the expression for $DFT\{v(n)\}$.





Transformation of Real-Valued Sequences – Part 8

DFT of a 2M-point real sequence – Part 3

The rearranging leads to

$$V_{2M}(\mu) = \sum_{n=0}^{M-1} \underbrace{v(2n)}_{=v_1(n)} W_{2M}^{2\mu n} + \sum_{n=0}^{M-1} \underbrace{v(2n+1)}_{=v_2(n)} W_{2M}^{(2n+1)\mu}$$
$$= \sum_{n=0}^{M-1} v_1(n) W_M^{\mu n} + W_{2M}^{\mu} \sum_{n=0}^{M-1} v_2(n) W_M^{n\mu}.$$

Finally we have:

$$V_{2M}(\mu) = V_{1,M}(\mu) + W_{2M}^{\mu} V_{2,M}(\mu), \quad \mu = 0, ..., 2M - 1.$$

Due to the Hermitian symmetry $V_{2M}(\mu) = V_{2M}^*(2M - \mu)$, μ only needs to be evaluated from 0 to M with $V_{i,M}(M) = V_{i,M}(0)$.





Transformation of Real-Valued Sequences – Part 9

DFT of a 2M-point real sequence – Part 4

Signal flow graph:



 \Rightarrow Computational savings by a factor of two compared to the complex-valued case since for real-valued input sequences only an *M*-point DFT is needed.



Transformation of Real-Valued Sequences – Part 10

DFT of a 2M-point real sequence – Part 5

Partner work – Please think about the following questions and try to find answers (first group discussions, afterwards broad discussion in the whole group).

□ How many memory elements do you need to store the result of a DFT of order M if the input sequence was real-valued?

□ Why is it useful to operate with complex-valued sequences?

General guestion) Where can you use a DFT/FFT? Application examples?



Christian-Albrechts-Universität zu Kiel

Summary

- □ Introduction
- Digital processing of continuous-time signals
- Efficient FIR structures
- □ DFT and FFT
 - □ DFT and signal processing
 - □ Fast computation of the DFT: The FFT
 - □ Transformation of real-valued sequences
 - □ Spectral and temporal refinement
- Digital filters
- Multi-rate digital signal processing





Spectral Refinement





Spectral Refinement

Basic Idea – Part 2

Basis setup

- Handsfree system with a DFT size of 256
- Hann window and with an overlap of 75 % (downsampling factor = 64)
- Processing power was about 30 % too high







Spectral Refinement

Basic Idea – Part 3

A standard DFT ...







Spectral Refinement

Basic Idea – Part 4





Spectral Refinement

Basic Idea – Part 5

Example









Spectral Refinement

Basic Idea – Part 6









Spectral Refinement

Basic Idea – Part 1









Spectral Refinement

Basic Idea – Part 7

Realization by FIR filters





CAU

Spectral Refinement

Different way to write a windowed DFT – Part 1

Input vector:

$$\boldsymbol{v}(nR) = \left[v(nR), v(nR-1), ..., v(nR-M+1) \right]^{\mathrm{T}}.$$

DFT (or FFT) of windowed input vector:

$$V(e^{j\Omega_{\mu}}, n) = \sum_{k=0}^{M-1} v(nR - k) h_k e^{-j\Omega_{\mu}k}$$

Equidistantly selected frequency supporting points:

$$\Omega_{\mu} = \frac{2\pi}{M}\mu$$
 with $\mu \in \{0, 1, ..., M-1\}.$



Spectral Refinement

Different way to write a windowed DFT – Part 2

A "windowed DFT" in matrix-vector notation:

$$\left[V(e^{j\Omega_0}, n), \, V(e^{j\Omega_1}, n), \, ..., \, V(e^{j\Omega_{M-1}}, n) \right]^{\mathrm{T}} = \boldsymbol{V}(n) = \boldsymbol{D}_M \, \boldsymbol{H} \, \boldsymbol{v}(nR).$$

DFT matrix:

$$\boldsymbol{D}_{M} = \begin{bmatrix} 1 & 1 & \cdots & 1 & 1 \\ 1 & e^{-j\frac{2\pi}{M}1} & \cdots & e^{-j\frac{2\pi}{M}1(M-2)} & e^{-j\frac{2\pi}{M}1(M-1)} \\ 1 & e^{-j\frac{2\pi}{M}2} & \cdots & e^{-j\frac{2\pi}{M}2(M-2)} & e^{-j\frac{2\pi}{M}2(M-1)} \\ \vdots & \vdots & \ddots & \vdots & & \vdots \\ 1 & e^{-j\frac{2\pi}{M}(M-1)} & \cdots & e^{-j\frac{2\pi}{M}(M-1)(M-2)} & e^{-j\frac{2\pi}{M}(M-1)(M-1)} \end{bmatrix}.$$

Window matrix:

 $oldsymbol{H} \,=\, \left[egin{array}{cccccccc} h_0 & 0 & \cdots & 0 & 0 \ 0 & h_1 & \cdots & 0 & 0 \ 0 & 0 & \cdots & 0 & 0 \ dots & dots & \ddots & dots & dots \ dots & dots & dots & dots & dots & dots \ dots & dots & dots & dots & dots \ dots & dots & dots & dots & dots \ dots & dots & dots & dots & dots \ dots & dots & dots & dots & dots \ dots & dots & dots & dots & dots \ dots & dots & dots & dots & dots \ dots & dots & dots & dots \ dots & dots & dots & dots \ dots & dots & dots \ dots$



Spectral Refinement

Basic Idea of spectral refinement – Part 1

Basic idea of spectral refinement:

$$\begin{bmatrix} V_{\text{ref}}(e^{j\Omega_0}, n) \\ V_{\text{ref}}(e^{j\Omega_1}, n) \\ \vdots \\ V_{\text{ref}}(e^{j\Omega_{M-1}}, n) \end{bmatrix} = V_{\text{ref}}(n) = S_{\text{ref}}\begin{bmatrix} V(n) \\ V(n-1) \\ \vdots \\ V(n-K+1) \end{bmatrix}$$





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Spectral Refinement

Basic Idea of spectral refinement – Part 2

Basic idea of spectral refinement:

$$\begin{bmatrix} V_{\rm ref}(e^{j\Omega_0}, n) \\ V_{\rm ref}(e^{j\Omega_1}, n) \\ \vdots \\ V_{\rm ref}(e^{j\Omega_{M-1}}, n) \\ V_{\rm int}(e^{j\Omega_0}, n) \\ V_{\rm int}(e^{j\Omega_1}, n) \\ \vdots \\ V_{\rm int}(e^{j\Omega_{M-1}}, n) \end{bmatrix} = \begin{bmatrix} \mathbf{V}_{\rm ref}(n) \\ \mathbf{V}_{\rm int}(n) \end{bmatrix} = \begin{bmatrix} \mathbf{S}_{\rm ref} \\ \mathbf{S}_{\rm int} \end{bmatrix} \begin{bmatrix} \mathbf{V}(n) \\ \mathbf{V}(n-1) \\ \vdots \\ \mathbf{V}(n-K+1) \end{bmatrix}.$$



Spectral Refinement

Basic Idea of spectral refinement – Part 3

Reordering (according to frequency) leads to :

$$\begin{bmatrix} V_{\mathrm{ref}}(e^{j\Omega_{0}}, n) \\ V_{\mathrm{int}}(e^{j\Omega_{0}}, n) \\ V_{\mathrm{ref}}(e^{j\Omega_{1}}, n) \\ V_{\mathrm{int}}(e^{j\Omega_{1}}, n) \\ \vdots \\ V_{\mathrm{ref}}(e^{j\Omega_{M-1}}, n) \\ V_{\mathrm{int}}(e^{j\Omega_{M-1}}, n) \end{bmatrix} = \begin{bmatrix} \widetilde{V}_{\mathrm{ref}}(e^{j\Omega_{0}}, n) \\ \widetilde{V}_{\mathrm{int}}(e^{j\Omega_{3}}, n) \\ \vdots \\ \widetilde{V}_{\mathrm{ref}}(e^{j\Omega_{M-1}}, n) \\ V_{\mathrm{int}}(e^{j\Omega_{M-1}}, n) \end{bmatrix} = \begin{bmatrix} \widetilde{V}_{\mathrm{ref}}(n) &= \underbrace{\mathbf{P} \begin{bmatrix} \mathbf{S}_{\mathrm{ref}} \\ \mathbf{S}_{\mathrm{int}} \end{bmatrix}}_{\widetilde{\mathbf{S}}} \begin{bmatrix} \mathbf{V}(n) \\ \mathbf{V}(n-1) \\ \vdots \\ \mathbf{V}(n-K+1) \end{bmatrix}.$$



Spectral Refinement

Basic Idea of spectral refinement – Part 4

The supervector of the DFT spectra can be reformulated:

$$\begin{bmatrix} \mathbf{V}(n) \\ \mathbf{V}(n-1) \\ \vdots \\ \mathbf{V}(n-K+1) \end{bmatrix} = \begin{bmatrix} \mathbf{D}_M \mathbf{H} \mathbf{v}(nR) \\ \mathbf{D}_M \mathbf{H} \mathbf{v}(n-1)R \\ \vdots \\ \mathbf{D}_M \mathbf{H} \mathbf{v}(n-K+1)R \end{bmatrix} = \begin{bmatrix} \mathbf{D}_M \mathbf{0}^{M \times M} \cdots \mathbf{0}^{M \times M} \\ \mathbf{0}^{M \times M} \mathbf{D}_M \cdots \mathbf{0}^{M \times M} \\ \vdots \\ \mathbf{0}^{M \times M} \mathbf{0}^{M \times M} \cdots \mathbf{D}_M \end{bmatrix} \begin{bmatrix} \mathbf{H} \mathbf{v}(nR) \\ \mathbf{H} \mathbf{v}(n-1)R \\ \vdots \\ \mathbf{H} \mathbf{v}(n-K+1)R \end{bmatrix}$$

Furthermore, we insert a longer time-domain vector:

$$\tilde{\boldsymbol{v}}(nR) = \left[v(nR), v(nR-1), ..., v(nR-M+1), ..., v(nR-M-(K-1)R+1]^{\mathrm{T}}, \right]$$



Spectral Refinement

Basic Idea of spectral refinement – Part 5

That allows as to rewrite the equation system as follows:

$$\begin{bmatrix} \mathbf{V}(n) \\ \mathbf{V}(n-1) \\ \vdots \\ \mathbf{V}(n-K+1) \end{bmatrix} = \begin{bmatrix} \mathbf{D}_{M} & \mathbf{0}^{M \times M} & \cdots & \mathbf{0}^{M \times M} \\ \mathbf{0}^{M \times M} & \mathbf{D}_{M} & \cdots & \mathbf{0}^{M \times M} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}^{M \times M} & \mathbf{0}^{M \times M} & \cdots & \mathbf{D}_{M} \end{bmatrix} \begin{bmatrix} \mathbf{H} \mathbf{v}(nR) \\ \mathbf{H} \mathbf{v}((n-1)R) \\ \vdots \\ \mathbf{H} \mathbf{v}(n-K+1)R \end{bmatrix}$$

$$= \begin{bmatrix} \mathbf{D}_{M} & \mathbf{0}^{M \times M} & \cdots & \mathbf{0}^{M \times M} \\ \mathbf{0}^{M \times M} & \mathbf{D}_{M} & \cdots & \mathbf{0}^{M \times M} \\ \mathbf{0}^{M \times M} & \mathbf{D}_{M} & \cdots & \mathbf{0}^{M \times M} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}^{M \times M} & \mathbf{0}^{M \times M} & \cdots & \mathbf{D}_{M} \end{bmatrix} \begin{bmatrix} \mathbf{H} & \mathbf{0}^{M \times R} & \cdots & \mathbf{0}^{M \times R} \\ \mathbf{0}^{M \times R} & \mathbf{H} & \cdots & \mathbf{0}^{M \times R} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}^{M \times R} & \mathbf{0}^{M \times R} & \cdots & \mathbf{H} \end{bmatrix} \tilde{\mathbf{v}}(nR).$$



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Spectral Refinement

Basic Idea of spectral refinement – Part 6

Combining everything lead to (for the refined and interpolated spectrum):

$$\widetilde{\boldsymbol{V}}_{\mathrm{ref}}(n) = \widetilde{\boldsymbol{S}}_{\mathrm{ref}} \begin{bmatrix} \boldsymbol{D}_{M} & \boldsymbol{0}^{M \times M} & \cdots & \boldsymbol{0}^{M \times M} \\ \boldsymbol{0}^{M \times M} & \boldsymbol{D}_{M} & \cdots & \boldsymbol{0}^{M \times M} \\ \vdots & \vdots & \ddots & \vdots \\ \boldsymbol{0}^{M \times M} & \boldsymbol{0}^{M \times M} & \cdots & \boldsymbol{D}_{M} \end{bmatrix} \begin{bmatrix} \boldsymbol{H} & \boldsymbol{0}^{M \times R} & \cdots & \boldsymbol{0}^{M \times R} \\ \boldsymbol{0}^{M \times R} & \boldsymbol{H} & \cdots & \boldsymbol{0}^{M \times R} \\ \vdots & \vdots & \ddots & \vdots \\ \boldsymbol{0}^{M \times R} & \boldsymbol{0}^{M \times R} & \cdots & \boldsymbol{H} \end{bmatrix} \widetilde{\boldsymbol{v}}(nR).$$



Spectral Refinement

A spectrum with higher resolution – Part 1

Starting again with a longer time-domain vector:

$$\tilde{\boldsymbol{v}}(nR) = \left[v(nR), v(nR-1), ..., v(nR-M+1), ..., v(nR-\widetilde{M}+1) \right]^{\mathrm{T}},$$

with $\widetilde{M} > M$ (e.g. $\widetilde{M} = 2M$).

And furthermore assuming that

$$\widetilde{M} = M + (K-1)R = 2M.$$

A second (high resolution) spectrum can be obtained by

$$\widetilde{\boldsymbol{V}}(n) = \boldsymbol{D}_{\widetilde{M}} \, \widetilde{\boldsymbol{H}} \, \widetilde{\boldsymbol{v}}(nR).$$



Spectral Refinement

A spectrum with higher resolution – Part 2

We will restrict the long window to be a sum of weighted and shifted short windows:

 $oldsymbol{A}\left[oldsymbol{h}^{\mathrm{T}},\,oldsymbol{h}^{\mathrm{T}},\,...,\,oldsymbol{h}^{\mathrm{T}}
ight] \;=\; \widetilde{oldsymbol{h}}.$

The weighting matrix is a sparse matrix of the form

$$oldsymbol{A} = ig[oldsymbol{A}_0, oldsymbol{A}_1, ..., oldsymbol{A}_{K-1}ig].$$

The first submatrix has the form

$$oldsymbol{A}_0 \;=\; \left[egin{array}{c} a_0 \, oldsymbol{I}^{M imes M} \ oldsymbol{0}^{(K-1)R imes M} \end{array}
ight]$$

The remaining submatrices are equivalently defined.





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Spectral Refinement

A spectrum with higher resolution – Part 3

Inserting these restrictions for the long window functions leads to: $\widetilde{V}(n) = D_{\widetilde{M}} \widetilde{H} \widetilde{v}(nR).$ $= D_{\widetilde{M}} A \begin{bmatrix} H & \mathbf{0}^{M \times R} & \cdots & \mathbf{0}^{M \times R} \\ \mathbf{0}^{M \times R} & H & \cdots & \mathbf{0}^{M \times R} \\ \vdots & \vdots & \vdots & \vdots & \vdots \end{bmatrix} \widetilde{v}(nR).$

$$\begin{bmatrix} \vdots & \vdots & \ddots & \vdots \\ \mathbf{0}^{M \times R} & \mathbf{0}^{M \times R} & \cdots & \mathbf{H} \end{bmatrix}$$

Comparing this with our refinement approach:

$$\widetilde{\boldsymbol{V}}_{\mathrm{ref}}(n) = \widetilde{\boldsymbol{S}}_{\mathrm{ref}} \begin{bmatrix} \boldsymbol{D}_{M} & \boldsymbol{0}^{M \times M} & \cdots & \boldsymbol{0}^{M \times M} \\ \boldsymbol{0}^{M \times M} & \boldsymbol{D}_{M} & \cdots & \boldsymbol{0}^{M \times M} \\ \vdots & \vdots & \ddots & \vdots \\ \boldsymbol{0}^{M \times M} & \boldsymbol{0}^{M \times M} & \cdots & \boldsymbol{D}_{M} \end{bmatrix} \begin{bmatrix} \boldsymbol{H} & \boldsymbol{0}^{M \times R} & \cdots & \boldsymbol{0}^{M \times R} \\ \boldsymbol{0}^{M \times R} & \boldsymbol{H} & \cdots & \boldsymbol{0}^{M \times R} \\ \vdots & \vdots & \ddots & \vdots \\ \boldsymbol{0}^{M \times R} & \boldsymbol{0}^{M \times R} & \cdots & \boldsymbol{H} \end{bmatrix} \widetilde{\boldsymbol{v}}(nR).$$





Spectral Refinement

A spectrum with higher resolution – Part 4

Setting

$$\widetilde{oldsymbol{V}}(n) \;=\; \widetilde{oldsymbol{V}}_{\mathrm{ref}}(n)$$

leads to

$$\widetilde{\boldsymbol{S}}_{\text{ref}} \begin{bmatrix} \boldsymbol{D}_{M} & \boldsymbol{0}^{M \times M} & \cdots & \boldsymbol{0}^{M \times M} \\ \boldsymbol{0}^{M \times M} & \boldsymbol{D}_{M} & \cdots & \boldsymbol{0}^{M \times M} \\ \vdots & \vdots & \ddots & \vdots \\ \boldsymbol{0}^{M \times M} & \boldsymbol{0}^{M \times M} & \cdots & \boldsymbol{D}_{M} \end{bmatrix} = \boldsymbol{D}_{\widetilde{M}} \boldsymbol{A}$$
$$\widetilde{\boldsymbol{S}}_{\text{ref}} = \boldsymbol{D}_{\widetilde{M}} \boldsymbol{A} \begin{bmatrix} \boldsymbol{D}_{M}^{-1} & \boldsymbol{0}^{M \times M} & \cdots & \boldsymbol{0}^{M \times M} \\ \boldsymbol{0}^{M \times M} & \boldsymbol{D}_{M}^{-1} & \cdots & \boldsymbol{0}^{M \times M} \\ \vdots & \vdots & \ddots & \vdots \\ \boldsymbol{0}^{M \times M} & \boldsymbol{0}^{M \times M} & \cdots & \boldsymbol{D}_{M}^{-1} \end{bmatrix}$$



Spectral Refinement





Spectral Refinement

Realization – Part 1

Example for a five tab refinement filter







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Spectral Refinement

Results – Part 1

Pitch estimation – data base

- Clean speech laryngograph/microphone database
- □ 38 speakers (18 female and 20 male)
- Mixing to different typical automotive SNR conditions

Pitch estimation – error types

- < 3 % absolute difference</p>
- < 10 % absolute difference</pre>
- < 20 % absolute difference</p>





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Spectral Refinement

Results – Part 2

Pitch estimation – results

 Correctness of estimated fundamental frequencies without and with spectral refinement for different SNR and tolerance ranges

		Accepted tolerance	High SNR	Medium SNR	Low SNR
nt	Standard method	< 3 %	62.1 %	70.1 %	47.2 %
		< 10 %	65.1 %	70.9 %	48.4 %
		< 20 %	65.5 %	71.4 %	49.3 %
	Spectral refinement up to 1 kHz	< 3 %	82.1 %	80.3 %	55.9 %
		< 10 %	88.1 %	85.3 %	58.2 %
		< 20 %	88.8 %	86.2 %	58.7 %
	Spectral refinement up to 3 kHz	< 3 %	83.2 %	80.1 %	53.4 %
		< 10 %	89.6 %	85.3 %	56.9 %
		< 20 %	90.6 %	86.3 %	57.5 %



Spectral Refinement

Outlook – Part 1

Polyphase filterbanks in the lecture "Adaptive Filters"





Spectral Refinement

Outlook – Part 2

Polyphase filterbanks in the lecture "Adaptive Filters"




DFT and FFT

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Summary

- □ Introduction
- Digital processing of continuous-time signals
- Efficient FIR structures
- □ DFT and FFT
 - DFT and signal processing
 - □ Fast computation of the DFT: The FFT
 - **Transformation of real-valued sequences**
 - Spectral refinement
- Digital filters
- Multi-rate digital signal processing



